

Segal Approach for Algebraic Structures

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Segal objects

Γ -spaces (Segal, 1974)

Denote by **Fin** the category of finite sets and by **Fin**₊ the category with the same objects as in **Fin**, and with morphisms given by partially defined maps, $S \supset U \rightarrow T$.

Denote 1 a one-element set. Given a finite set S , its elements $\{s\} \subset S$ induce the morphisms $\rho_s : S \leftarrow 1 \rightarrow 1$ in **Fin**₊. A Γ -space is then a functor

$$A : \mathbf{Fin}_+ \longrightarrow \mathbf{Top},$$

such that the induced morphism (*the Segal map*)

$$\prod A(\rho_S) : A(S) \longrightarrow A(1)^S$$

is a (weak) homotopy equivalence for each $S \in \mathbf{Fin}_+$.

Homotopy coherent multiplication

Given $S \in \mathbf{Fin}_+$, we have the usual map $\pi_S : S \xrightarrow{=} S \rightarrow 1$ defined everywhere on S . Consider the span

$$\begin{array}{ccc} & A(S) & \\ \prod A(\rho_S) \swarrow & & \searrow A(\pi_S) \\ & A(1)^S & A(1). \end{array} \quad \sim$$

which left arrow is a homotopy equivalence. Inverting it, we obtain multiplication operations $m_S : A(1)^S \rightarrow A(1)$.

In $\mathbf{Ho Top}$, the homotopy type of $A(1)$ is a commutative monoid. But a Γ -space contains much more information than an H -space (“Segal delooping machinery”).

Operator categories (Barwick 2013)

One can replace **Fin** by **Ord**, the category of finite ordered sets. Somewhat equivalently, one can work with Δ^{op} instead of **Fin**₊. A way to systematise:

Definition. An *operator category* \mathbf{C} is a small category with a terminal object 1 , such that the hom-sets $\mathbf{C}(1, x)$ are finite for each $x \in \mathbf{C}$ and that pullbacks exist along any map $1 \rightarrow x$ of \mathbf{C} .

Examples. Already mentioned **Fin**, **Ord**. Consider also the category **B**:

1. Its objects are injections $f : S \hookrightarrow D$ with domain $S \in \mathbf{Fin}$, into the 2-disc D same thing as $|S|$ distinct points in D .
2. A map between $f : S \hookrightarrow D$ and $f' : S' \hookrightarrow D$ is given by a set map $\alpha : S \rightarrow S'$ and by a path f to $f' \circ \alpha$ in the stratified groupoid $\Pi_1^{EP}(Cf(S, D))$ of the configuration space $Cf(S, D) = \{S \rightarrow D\}$.

Intuition: **B** is like **Fin**, but with braid-automorphisms.

Algebra Classifiers

Definition. Let \mathbf{C} be an operator category. Its *algebra classifier* is the category \mathbf{C}_+ such that

1. $\text{Ob } \mathbf{C}_+ = \text{Ob } \mathbf{C}$,
2. the hom-sets $\mathbf{C}_+(x, y)$ are given by equivalence classes of span diagrams $x \leftarrow z \rightarrow y$, where $z \rightarrow y$ is in \mathbf{C} and $z \hookrightarrow x$ is an *admissible monomorphism*: a composition of pullbacks of elementary (admissible) monos $1 \rightarrow t$.

For example, \mathbf{Fin}_+ is the same category as before. For \mathbf{Ord} , the admissible monos are interval inclusions. All monos are admissible in \mathbf{B} .

A morphism in \mathbf{C}_+ is *inert* if it can be presented as $z \leftarrow x \xrightarrow{=} x$, and *active* if it can be presented as $y \xleftarrow{=} y \rightarrow t$. Inert and active maps form a factorisation system $(\text{In}_{\mathbf{C}}, \text{Act}_{\mathbf{C}})$ on \mathbf{C}_+ .

Homotopical algebra with Segal objects

Given an operator category \mathbf{C} , we can now define a Segal \mathbf{C} -spaces as a functor $A : \mathbf{C}_+ \rightarrow \mathbf{Top}$ with Segal conditions: that

$$\prod A(\rho_x) : A(x) \longrightarrow A(1)^{|x|}$$

is a homotopy equivalence for $x \in \mathbf{C}_+$, $|x| = \mathbf{C}(1, x)$.

For example, Segal **Ord**-spaces are homotopy associative monoids.

Because of the correspondence between constructible sheaves on $Cf(S, D)$ and functors $\Pi_1^{EP}(Cf(S, D)) \rightarrow \mathbf{Top}$, Segal **B**-spaces are the same data as constructible factorisation spaces over a 2-disc, hence E_2 -algebras in \mathbf{Top} .

Barwick shows how to construct operator categories related to E_n -spaces. There are also other approaches (Batanin-Markl, Berger).

Resolutions

Resolutions (of operator categories)

Definition. A functor $F : \mathcal{D} \rightarrow \mathcal{C}$ is a *resolution* if for each $\mathbf{c}_{[n]} = c_0 \rightarrow \dots \rightarrow c_n$ de \mathbb{C} , the category $\mathcal{D}(\mathbf{c}_{[n]})$ of strings $d_0 \rightarrow \dots \rightarrow d_n$ in $\text{Fun}([n], \mathcal{D})$ with isomorphisms $(Fd_0 \rightarrow \dots \rightarrow Fd_n) \cong \mathbf{c}_{[n]}$ is contractible.

Formal examples: $\int N\mathcal{C} \rightarrow \mathcal{C}^{\text{op}}$, $\int N\mathcal{C} \rightarrow \mathcal{C}$, equivalences of categories, (op)fibrations with contractible fibres.

Definition. A functor $F : \mathcal{D} \rightarrow \mathcal{C}$ between operator categories is a resolution if:

1. The functor F preserves limits and $\mathcal{D}(1, x) \cong \mathcal{C}(1, F(x))$,
2. The functor F is a resolution.

NB: No requirement for $\mathcal{D}_+ \rightarrow \mathcal{C}_+$ to be a resolution!

Example: classifying spaces and representations

Let G and denote BG its classifying space. Let I be a regular cell decomposition of BG , viewed as a poset. There is a natural functor

$$F : I \longrightarrow \Pi_1(BG) \cong G$$

given by choosing a point in each cell. The functor F is a resolution.

Denoting $\mathcal{D}(\mathcal{C}, k)$ the derived category of functors $\mathcal{C} \rightarrow \mathbf{DVect}_k$, we have

$$F^* : \mathbf{DRep}(G, k) \cong \mathcal{D}(\Pi_1(BG), k) \rightarrow \mathcal{D}(I, k) \cong \mathcal{D}(A(I)\text{-}\mathbf{Mod}).$$

One can prove that F^* is full and faithful, and that its image consists of $X : I \rightarrow \mathbf{DVect}_k$ which are *locally constant*: for each $f : i \rightarrow i'$ de I , $X(f)$ is a quasi-isomorphism.

Does this example arise in operator categories?

Planar trees

A *planar tree* T is

1. a connected unoriented graph without cycles possessing a distinguished vertex of valency 1 called the root,
2. both the set of vertices $V(T)$ and the set of edges $E(T)$ are finite,
3. for each $v \in V(T)$ there is a datum of cyclic order on the set of edges attached to v . This makes T into an oriented graph.

A morphism $f : T \rightarrow T'$ is an oriented cellular map $|f| : |T| \rightarrow |T'|$ between the geometric realisations, such that

1. $|f|$ preserves the roots,
2. for each $a, b \in V(T)$, the image by $|f|$ of a geodesic between a and b in $|T|$ is a geodesic between $|f|(a)$ et $|f|(b)$.

Planar trees form a category \mathbf{T}_0 . It is an operator category with zero object.

Stable planar trees

A *marked* planar tree is a pair (T, S) , where $T \in \mathbf{T}_0$ and $S \subset V(T)$ is a subset not containing the root. A marked planar tree is *stable* if each non-marked vertex (but the root) has valency at least three.

Definition. An object of the category \mathbf{T} is a marked stable planar tree (T, S) . A morphism $(T, S) \rightarrow (T', S')$ is given by a map $f : T \rightarrow T'$ in \mathbf{T}_0 such that f sends S to S' .

The category \mathbf{T} is an operator category.

There is another category $\tilde{\mathbf{T}}$ with the objects given by those of \mathbf{T} plus an immersion into a 2-disc which sends all roots to one fixed point on the boundary. The forgetful functor $\tilde{\mathbf{T}} \rightarrow \mathbf{T}$ is an equivalence of categories.

Forgetting everything but the marked vertices induces another functor $\tilde{\mathbf{T}} \rightarrow \mathbf{B}$. Inverting the equivalence $\tilde{\mathbf{T}} \xrightarrow{\sim} \mathbf{T}$, we get $F : \mathbf{T} \rightarrow \mathbf{B}$.

Resolution of \mathbf{B} by planar trees

Theorem PT (partially observed by Kontsevich-Soibelman, Kaledin).

The functor $F : \mathbf{T} \rightarrow \mathbf{B}$ is a resolution of operator categories.

Corollary. *The inverse image functor between categories of Segal objects*

$$F^* : \mathrm{Ho} \mathrm{Fun}_{\mathrm{Seg}}(\mathbf{B}_+, \mathbf{Top}) \rightarrow \mathrm{Ho} \mathrm{Fun}_{\mathrm{Seg}}(\mathbf{T}_+, \mathbf{Top})$$

is full and faithful, and its essential image consists of $A : \mathbf{T}_+ \rightarrow \mathbf{Top}$ such that whenever $F(f)$ is an iso in \mathbf{B}_+ , the image $A(f)$ is a weak homotopy equivalence.

Thus to construct E_2 -algebras in \mathbf{Top} , do it first over \mathbf{T} in a good manner, then descend. Examples might involve another proof of

Deligne Conjecture. *For a dg-algebra A , there is an E_2 -algebra structure on $CH^\bullet(A, A)$ (which realises the well-known operations on $HH^\bullet(A, A)$).*

...except how to do Segal objects in non-cartesian monoidal setting?

Grothendieck (op)fibrations

Opcartesian morphisms (à l'ancienne)

Let $p : \mathcal{E} \rightarrow \mathcal{C}$ be a functor. A morphism $\alpha : x \rightarrow y$ of \mathcal{E} is *p-opcartesian* if for each $\beta : x \rightarrow z$ such that $p\beta = p\alpha$ there exists unique factorisation $\beta = \gamma\alpha$, where $p(\gamma) = id_{p(y)}$:

$$\begin{array}{ccc} & & z \\ & \nearrow \beta & \uparrow \exists! \gamma \\ x & \xrightarrow{\alpha} & y \\ \downarrow & & \downarrow \\ px & \xrightarrow{p\alpha} & py \end{array}$$

This definition is from SGA1; modern references call this notion locally opcartesian.

Opfibrations

A functor $p : \mathcal{E} \rightarrow \mathcal{C}$ is a Grothendieck opfibration if:

1. For each $f : c \rightarrow c'$ of \mathcal{C} and $x \in \mathcal{E}$ with $px = c$ there exists an opcartesian lifting $\alpha : x \rightarrow f!x$, $p\alpha = f$:

$$\begin{array}{ccc} x & \overset{\exists \alpha}{\dashrightarrow} & f!x \\ \downarrow & & \downarrow \\ c & \xrightarrow{f} & c'. \end{array}$$

2. The composition of opcartesian morphisms is opcartesian.

For $c \in \mathcal{C}$, denote by $\mathcal{E}(c) = p^{-1}c$ the fibre of p over c . A choice of opcartesian liftings along $f : c \rightarrow c'$ defines a functor $f! : \mathcal{E}(c) \rightarrow \mathcal{E}(c')$.

Dual notions: cartesian maps, Grothendieck fibrations.

Example: symmetric monoidal categories (Lurie, Segal)

Given a symmetric monoidal category \mathcal{M} with \otimes , construct an opfibration $\mathcal{M}^{\otimes} \rightarrow \mathbf{Fin}_+$ as follows.

1. An object of \mathcal{M}^{\otimes} is a pair consisting of $S \in \mathbf{Fin}_+$ and a S -indexed family $\{X_s\}_{s \in S}$ of objects of \mathcal{M} .
2. A morphism in \mathcal{M}^{\otimes} , $(S, \{X_s\}_{s \in S}) \rightarrow (T, \{Y_t\}_{t \in T})$, consists of a map $f : S \rightarrow T$ in \mathbf{Fin}_+ together with maps $\otimes_{s \in f^{-1}(t)} X_s \rightarrow Y_t$ for each $t \in T$.
3. The projection $(S, \{X_s\}_{s \in S}) \rightarrow S$ gives a functor $\mathcal{M}^{\otimes} \rightarrow \mathbf{Fin}_+$.

We see that $\mathcal{M}^{\otimes}(S) \cong \mathcal{M}^S$. The map $S \mapsto \mathcal{M}^{\otimes}(S)$ can be made into a (pseudo-)functor which satisfies Segal conditions in \mathbf{Cat} . For other monoidal structures (associative, braided) one can make similar considerations.

Algebras as sections

Let $p : \mathcal{E} \rightarrow \mathcal{C}$ be an opfibration. A *section* of p is a functor $A : \mathcal{C} \rightarrow \mathcal{E}$ such that $pA = id$. The sections of p form a category $\text{Sect}(\mathcal{C}, \mathcal{E})$ with fibrewise natural transformations.

In the example of $\mathcal{M}^{\otimes} \rightarrow \mathbf{Fin}_+$, consider a section $A : \mathbf{Fin}_+ \rightarrow \mathcal{M}^{\otimes}$ such that for each inert morphism

$$j : S \longleftarrow \supset T \xrightarrow{=} T,$$

the morphism $A(j)$ is opcartesian. Then we get $A(S) \cong (A(1), \dots, A(1))$ and we obtain maps $A(1)^{\otimes S} \rightarrow A(1)$ in $\mathcal{M}^{\otimes}(1) = \mathcal{M}$. This way, the object $A(1)$ becomes a commutative monoid in \mathcal{M} , and this construction can be reversed.

However, commutative algebras in $\mathcal{M} = \mathbf{DVect}_k$ do not form a good homotopical category.

Homotopy theory of sections?

One way to remedy the issue consists of doing $\mathcal{M}^{\otimes} \rightarrow \mathbf{Fin}_+$ and similarly for more general opfibrations $\mathcal{E} \rightarrow \mathcal{C}$ is to pass to higher-categorical context, taking a higher localisation of \mathcal{E} . Associated difficulties arise (“very fibrant replacement”).

We still have no Segal description for monoids in \mathcal{M} . In general, an ordinary section A of an opfibration $\mathcal{E} \rightarrow \mathcal{C}$ produces, out of $f : c \rightarrow c'$, a map $f_! A(c) \rightarrow A(c')$. Can we get a “weak section” X , which, out of $f : c \rightarrow c'$, would produce a diagram

$$\begin{array}{ccc} & X(f) & \\ & \swarrow & \searrow \\ f_! X(c) & \sim & X(c') \end{array}$$

with left arrow a weak equivalence?

The approach proposed below addresses the latter point, and works in categorical or higher-categorical setting.

Derived, or Segal, sections

Simplicial replacements

For a small category \mathcal{C} , its *simplicial replacement* is the category \mathbb{C} such that

1. An object $\mathbf{c}_{[n]} \in \mathbb{C}$ is a sequence of composable arrows of \mathcal{C} :

$$\mathbf{c}_{[n]} = c_0 \rightarrow c_1 \rightarrow \dots \rightarrow c_n.$$

2. A morphism $\alpha : \mathbf{c}_{[n]} \rightarrow \mathbf{c}'_{[m]}$ is given by $a : [m] \rightarrow [n]$ in Δ such that $c_{a(i)} = c'_i$ for each $i \in [m]$.

Defined this way, $\mathbb{C} = (\int N\mathcal{C})^{\text{op}}$. The maps $\mathbf{c}_{[n]} \mapsto c_0$, $\mathbf{c}_{[n]} \mapsto c_n$ yield functors $\mathbb{C} \xrightarrow{h} \mathcal{C}$ and $\mathbb{C} \xrightarrow{t} \mathcal{C}^{\text{op}}$.

Extension of $\mathcal{E} \rightarrow \mathcal{C}$ to \mathbb{C}

Recall the final object map $t : \mathbb{C} \rightarrow \mathcal{C}^{\text{op}}$, $\mathbf{c}_{[n]} \mapsto c_n$. We want to use it to lift the opfibration (covariant family) $\mathcal{E} \rightarrow \mathcal{C}$ to \mathbb{C} . However, for this we have to replace $\mathcal{E} \rightarrow \mathcal{C}$ by its *transposed fibration* (contravariant family) $\mathcal{E}^\top \rightarrow \mathcal{C}^{\text{op}}$.

This family is characterised by the facts that $\mathcal{E}^\top(c) \cong \mathcal{E}(c)$ and that for each morphism $f : c' \leftarrow c$ of \mathcal{C}^{op} , the transition functor $\mathcal{E}^\top(c) \rightarrow \mathcal{E}^\top(c')$ is isomorphic to $f_! : \mathcal{E}(c) \rightarrow \mathcal{E}(c')$. This is, however, a fibration, so a normed section of $(\mathcal{M}^\otimes)^\top \rightarrow \mathbf{Fin}_+^{\text{op}}$ would correspond to a *coalgebra* object in \mathcal{M} .

This is natural for Segal formalism, which treats algebraic objects as “bigger” coalgebraic objects with special conditions.

We can now consider $t^*\mathcal{E}^\top \rightarrow \mathbb{C}$, the inverse image of $\mathcal{E}^\top \rightarrow \mathcal{C}^{\text{op}}$ along $t : \mathbb{C} \rightarrow \mathcal{C}^{\text{op}}$.

Sections of $t^* \mathcal{E}^\top \rightarrow \mathbb{C}$

Let us consider a section $X : \mathbb{C} \rightarrow t^* \mathcal{E}^\top$ of the fibration $t^* \mathcal{E}^\top \rightarrow \mathbb{C}$. For $f : c \rightarrow c'$ (associated functor $f_! : \mathcal{E}(c) \rightarrow \mathcal{E}(c')$), we get a span in $\mathcal{E}(c')$ of the desired form:

$$\begin{array}{ccc} c & \xrightarrow{f} & c' \\ & \swarrow & \searrow \\ & c & c' \end{array} \quad \mapsto \quad \begin{array}{ccc} & X(c \rightarrow c') & \\ & \swarrow & \searrow \\ f_! X(c) & & X(c'). \end{array}$$

If $\mathcal{E} \rightarrow \mathbb{C}$ is equipped with a homotopical structure (e.g. weak equivalences in each $\mathcal{E}(c)$ preserved by $f_!$), we can ask for the left arrow to be a weak equivalence.

This is a naïve picture to keep, and works well in higher-categorical setting. However, in model-categorical setting, such objects are insufficiently homotopy coherent.

Simplicial extension of an opfibration

Given $\mathcal{E} \rightarrow \mathcal{C}$, its *simplicial extension* $\mathbf{E} \rightarrow \mathbb{C}$ is a family with fibres given by $\mathbf{E}(\mathbf{c}_{[n]}) = \text{Sect}([n]^{\text{op}}, \mathbf{c}_{[n]}^* \mathcal{E}^\top)$, where we consider $\mathbf{c}_{[n]}$ as a functor:

$$\begin{array}{ccc} & & \mathcal{E}^\top \\ & \nearrow \text{dashed} & \downarrow \\ [n]^{\text{op}} & \xrightarrow{\mathbf{c}_{[n]}} & \mathcal{C}^{\text{op}} \end{array}$$

If the opfibration $\mathcal{E} \rightarrow \mathcal{C}$ is fibrewise (finitely) complete, then the family $\mathbf{E} \rightarrow \mathbb{C}$ is a bifibration (trivially an opfibration, but also a fibration).

Definition. The category of *presections* is the category $\text{PSect}(\mathcal{C}, \mathcal{E}) := \text{Sect}(\mathbb{C}, \mathbf{E})$. A presection $X : \mathbb{C} \rightarrow \mathbf{E}$ is derived (or Segal) if the image $X(\alpha)$ of any left interval inclusion $\alpha : \mathbf{c}_{[n]} \rightarrow \mathbf{c}'_{[m]}$ factors as a weak equivalence followed by cartesian. We thus have

$$\text{DSect}(\mathcal{C}, \mathcal{E}) \subset \text{PSect}(\mathcal{C}, \mathcal{E}) = \text{Sect}(\mathbb{C}, \mathbf{E}).$$

The model structure

Model category \mathbf{PSect}

Let $\mathcal{E} \rightarrow \mathcal{C}$ be a *model opfibration*, that is, each $\mathcal{E}(x)$ is a model category, and the transition functors $\mathcal{E}(x) \rightarrow \mathcal{E}(y)$ preserve weak equivalences and fibrations. (Think $\mathbf{DVect}_k^{\otimes} \rightarrow \mathbf{Fin}_+$)

Theorem. *In this case, the presections category $\mathbf{PSect}(\mathcal{C}, \mathcal{E}) = \mathbf{Sect}(\mathcal{C}, \mathbf{E})$ possesses a model structure, with weak equivalences fibrewise.*

Implication: the category $\mathbf{DSect}(\mathcal{C}, \mathcal{E})$ is realised as a full homotopical subcategory of a model category $\mathbf{PSect}(\mathcal{C}, \mathcal{E})$. Denote by $\mathbf{Ho PSect}(\mathcal{C}, \mathcal{E})$ and $\mathbf{Ho DSect}(\mathcal{C}, \mathcal{E})$ the corresponding localisations.

This result is a consequence of a more general theorem for families of model categories over Reedy categories.

Semifibrations

A *semifibration* over a factorisation category $(\mathcal{C}, \mathcal{L}, \mathcal{R})$ is a functor $p : \mathcal{S} \rightarrow \mathcal{C}$ such that

1. for each morphism $l : x \rightarrow y$ of \mathcal{L} and $Y \in \mathcal{S}(y)$ there is a cartesian lift $l^* Y \rightarrow Y$ of l ,
2. for each morphism $r : x \rightarrow y$ of \mathcal{R} and $X \in \mathcal{S}(x)$ there is an opcartesian lift $X \rightarrow \eta X$ of r ,
3. given a morphism $\alpha : X \rightarrow Y$ of \mathcal{S} and a factorisation of $p(\alpha)$ as $x \xrightarrow{r} z \xrightarrow{l} y$ (wrong arrow order), there is a decomposition of α as

$$X \xrightarrow{\rho} Z \xrightarrow{\omega} Z' \xrightarrow{\lambda} Y,$$

such that $p(\rho) = r$, $p(\lambda) = l$ and $p(\omega) = id_z$.

Theorem MS

Let \mathcal{R} be a Reedy category. A *model semifibration* over \mathcal{R} is a semifibration $\mathcal{S} \rightarrow \mathcal{R}$ for the Reedy factorisation system $(\mathcal{R}, \mathcal{R}_-, \mathcal{R}_+)$ such that

1. each fibre $\mathcal{S}(x)$ is a model category,

2L. for each $l : x \rightarrow y$ de \mathcal{R}_- , the transition functor $l^* : \mathcal{S}(y) \rightarrow \mathcal{S}(x)$ preserves fibrations and trivial fibrations,

3L. for each x in \mathcal{R} , either

the matching category $Mat(x)$ is a disjoint union of categories with initial objects, or

the functor $\text{Sect}(Mat(x), \mathcal{S}) \rightarrow Fun(Mat(x), \mathcal{S}(x))$ preserves limits,

and dually, **2R, 3R.**

Theorem MS. *The category $\text{Sect}(\mathcal{R}, \mathcal{S})$ of sections of a model semifibration $\mathcal{S} \rightarrow \mathcal{R}$ has a model structure, in which weak equivalences are fibrewise, and the fibrations and cofibrations are Reedy.*

Theorem MS: discussion

When $\mathcal{S} \rightarrow \mathcal{R}$ is a bifibration, the result reduces to that of Hirschowitz-Simpson (theory of Quillen presheaves).

Corollary. Let $\mathcal{E} \rightarrow \mathcal{C}$ be a model opfibration, then the simplicial extension $\mathbf{E} \rightarrow \mathbb{C}$ is a model semifibration.

Proof. Each fibre $\mathbf{E}(\mathbf{c}_{[n]}) = \text{Sect}([n]^{\text{op}}, \mathcal{E}^\top)$ is a model category by Theorem MS, and then we apply Theorem MS (or H.-S.) again, globally to $\mathbf{E} \rightarrow \mathbb{C}$.

Contrary to H.-S., we have a case in which nothing is assumed on transition functors (adjoints, exactness...). This allows us to consider n -fold tensor products.

For a fibrewise-presentable, accessible higher opfibration $\mathcal{E} \rightarrow \mathcal{C}$ over a 1-category, the presentability of $\text{PSect}(\mathcal{C}, \mathcal{E}) = \text{Sect}(\mathbb{C}, \mathbf{E})$ is almost readily apparent.

Resolutions and Segal sections

Locally constant derived sections

Let $\mathcal{E} \rightarrow \mathcal{C}$ be a model opfibration and $Iso(\mathcal{C}) \subset \mathcal{S} \subset \mathcal{C}$ a subcategory.

Définition. A derived section $X \in \text{DSect}(\mathcal{C}, \mathcal{E})$ is \mathcal{S} -*locally constant* if X sends to weakly cartesian arrows those maps $\mathbf{c}_{[n]} \rightarrow \mathbf{c}'_{[m]}$ which verify the following

1. the induced morphism $[m] \rightarrow [n]$ in Δ is a right interval inclusion,
2. the maps $c_{i-1} \rightarrow c_i$, $1 \leq i \leq n-1$, belong to \mathcal{S} .

Example. Each algebra $A : \mathbf{Fin}_+ \rightarrow \mathcal{M}^\otimes$ gives a derived section locally constant along the inert morphisms $In_{\mathbf{Fin}}$.

Denote by $\text{DSect}_{\mathcal{S}}(\mathcal{C}, \mathcal{E}) \subset \text{DSect}(\mathcal{C}, \mathcal{E})$ the subcategory of derived \mathcal{S} -locally constant sections.

Theorem RES

Let $\mathcal{E} \rightarrow \mathcal{C}$ be a model opfibration, $Iso(\mathcal{C}) \subset \mathcal{S} \subset \mathcal{C}$ a subcategory and $F : \mathcal{D} \rightarrow \mathcal{C}$ a functor. Then the functor F induces

$$F^* : \text{DSect}_{\mathcal{S}}(\mathcal{C}, \mathcal{E}) \longrightarrow \text{DSect}_{F^*\mathcal{S}}(\mathcal{D}, F^*\mathcal{E}),$$

where $F^*\mathcal{E} \rightarrow \mathcal{D}$ is the pullback of $\mathcal{E} \rightarrow \mathcal{C}$, and $F^*\mathcal{S} \subset \mathcal{D}$ is a subcategory given by those f of \mathcal{D} such that $F(f) \in \mathcal{S}$.

Theorem RES. *If moreover $F : \mathcal{D} \rightarrow \mathcal{C}$ is a resolution, then*

$$\text{h}F^* : \text{Ho DSect}_{\mathcal{S}}(\mathcal{C}, \mathcal{E}) \longrightarrow \text{Ho DSect}_{F^*\mathcal{S}}(\mathcal{D}, F^*\mathcal{E})$$

is an equivalence of categories.

If $\mathcal{S} = Iso(\mathcal{C})$, then $F^*Iso(\mathcal{C})$ is a subcategory of morphisms of \mathcal{D} which become isomorphisms in \mathcal{C} .

Some comments on the proof

To prove Theorem RES, we construct a functor $hF_!$ inverse to hF^* , as

$$hF_! := \mathbb{L}p_{F,!} \circ h\mu_F^* \circ \mathbb{R}\delta_{\mathcal{D},*}, \text{ where:}$$

1. The functor $\delta_{\mathcal{D},*} : \text{PSect}(\mathcal{D}, \mathcal{E}) = \text{Sect}(\mathbb{D}, \mathbf{E}) \rightarrow \text{Sect}(\mathbb{D}_\Pi, \mathbf{E}_\Pi)$ is a right Kan extension along $\delta_{\mathcal{D}} : \mathbb{D} \rightarrow \mathbb{D}_\Pi$. The category \mathbb{D}_Π is the Π -replacement of \mathcal{D} , its objects are $\mathbf{c}_P : P \rightarrow \mathcal{D}$, with $P \in \Pi$ a finite poset with initial and final objects.
2. The functor $\mu_F^* : \text{Sect}(\mathbb{D}_\Pi, \mathbf{E}_\Pi) \rightarrow \text{Sect}(\mathbb{T}(F), \mu^*\mathbf{E}_\Pi)$ is the inverse image along $\mu : \mathbb{T}(F) \rightarrow \mathbb{D}_\Pi$. Here, we note by $p_F : \mathbb{T}(F) \rightarrow \mathbb{C}$ the tower of F , an opfibration which fibres are simplicial replacements of $\mathcal{D}(\mathbf{c}_{[n]})$.
3. The functor

$$p_{F,!} : \text{Sect}(\mathbb{T}(F), \mu^*\mathbf{E}_\Pi) \rightarrow \text{Sect}(\mathbb{T}(F), p_F^*\mathbf{E}) \rightarrow \text{Sect}(\mathbb{C}, \mathbf{E}) = \text{PSect}(\mathbb{C}, \mathcal{E})$$

is obtained from a left Kan extension along the opfibration $p_F : \mathbb{T}(F) \rightarrow \mathbb{C}$.

Derived algebras

Recall: For an operator category C , the algebra classifier C_+ consists of partially defined maps with admissible domain.

Definition. Let C be an operator category. An C -monoidal category is a Grothendieck opfibration $\mathcal{M}^\otimes \rightarrow C_+$ such that for each $x \in C_+$, the induced functor

$$\mathcal{M}^\otimes(x) \longrightarrow \prod_{(x \rightarrow 1) \in \text{In}_C} \mathcal{M}^\otimes(1)$$

is an equivalence of categories.

A C -monoidal model category is a C -monoidal C category $\mathcal{M}^\otimes \rightarrow C_+$ which is also a model opfibration. (use suitable presentability/accessibility for highercat setting)

Definition. Given an C -monoidal model category, its *category of derived algebras* is $\text{DAlg}(C, \mathcal{M}) := \text{DSect}_{\text{In}_C}(C_+, \mathcal{M}^\otimes)$, that is the category of derived sections of $\mathcal{M}^\otimes \rightarrow C_+$ which are In_C -locally constant.

Theorem RES-ALG

Definition (reminder). A functor $F : \mathbf{D} \rightarrow \mathbf{C}$ between operator categories is a resolution if:

1. The functor F preserves limits and $\mathbf{D}(1, x) \cong \mathbf{C}(1, F(x))$,
2. The functor F is a resolution.

Theorem RES-ALG. *Given a \mathbf{C} -monoidal model category $\mathcal{M}^{\otimes} \rightarrow \mathbf{C}_+$ and a resolution of operator categories $F : \mathbf{D} \rightarrow \mathbf{C}$, the induced functor*

$$hF^* : \mathrm{Ho} \mathrm{DAlg}(\mathbf{C}, \mathcal{M}) \longrightarrow \mathrm{Ho} \mathrm{DAlg}_{F^* \mathrm{Iso}(\mathbf{C})}(\mathbf{D}, F^* \mathcal{M})$$

is an equivalence of categories, where $\mathrm{DAlg}_{F^ \mathrm{Iso}(\mathbf{C})}(\mathbf{D}, F^* \mathcal{M})$ is the category of derived algebras locally constant along $F^* \mathrm{Iso}(\mathbf{C}) \subset \mathbf{D} \subset \mathbf{D}_+$.*

Preuve. Repeated “black box” application of Theorem RES.

Resolution of \mathbf{B} and Segal algebras

Theorem PT (reminder). *There is a functor $F : \mathbf{T} \rightarrow \mathbf{B}$ which is a resolution of operator categories.*

Theorems PT and RES-ALG imply that the inverse image functor

$$hF^* : \mathrm{Ho} \mathrm{DAlg}(\mathbf{B}, \mathcal{M}) \rightarrow \mathrm{Ho} \mathrm{DAlg}_{F^* \mathrm{Iso}(\mathbf{B})}(\mathbf{T}, \mathcal{M})$$

is an equivalence of categories.

This can be used to prove the Deligne conjecture outside of the operad formalism. For $\mathbf{DVect}_k^\otimes \rightarrow \Gamma_+$ and a dg -algebra A over k , there is a combinatorial way to construct a derived algebra $CH_{\mathbf{T}}^\bullet(A) \in \mathrm{DAlg}(\mathbf{T}, \mathbf{DVect}_k)$ whose value at $1 \in \mathbf{T}$ is $CH^\bullet(A, A)$ and which is locally constant.

Sketch of construction

Over \mathbf{T}_+ , there is an opfibration $\mathbf{Bimod}_A^{\mathbf{T}} \rightarrow \mathbf{T}_+$ with fibres over (T, S) equivalent to $\prod_{v \in S} (A^{\otimes out(v)} \otimes A\text{-Bimod})$ (bimodules viewed as functors of many arguments).

This opfibration has two distinguished sections $L(A), R(A)$, induced by the bimodules $A^{\otimes out(v)} \otimes A$ and $\text{Hom}_k(A^{\otimes out(v)}, A)$ in each fibre, respectively.

Taking a hom-pairing between the corresponding derived sections (amounts to projectively deriving $L(A)$) produces $CH_{\mathbf{T}}^{\bullet}(A) \in \text{DAlg}(\mathbf{T}, \mathbf{DVect}_k)$.

Descending the obtained derived section to \mathbf{B} gives us

$$CH_{\mathbf{B}}^{\bullet}(A) \in \text{DAlg}(\mathbf{B}, \mathbf{DVect}_k),$$

a presentation of $CH^{\bullet}(A, A)$ as an E_2 -algebra.

Thank you.